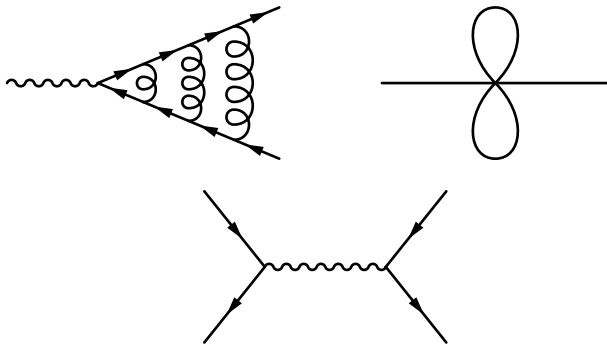


Hopf Algebras & Feynman Graphs

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PART I: COMBINATORICS

- Rooted Trees
- Connes–Kreimer Hopf Algebra of Rooted Trees

PART II: COMBINATORIAL PHYSICS

- a hint of Physics
- Feynman Graphs
- Combinatorial Physical Theories
- Renormalization Hopf Algebras
- Rooted Trees & Feynman Graphs

Definition

A **combinatorial class** is a set C equipped with a size function $\|\cdot\| : C \rightarrow \mathbb{Z}_{n \geq 0}$, such that $C_n := \{c \in C \mid |c| = n\}$ is finite $\forall n$.

- C : *configuration space* (e.g., finite graphs)
- C_n : *macrostate* (e.g., finite graphs with 4 vertices)
- $c \in C_n$: *microstate* (e.g., chain graph of 4 vertices)

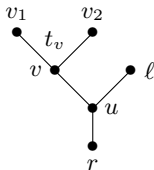
Important: combinatorial class T of *planar rooted trees*
 $|\text{tree}| = \#(\text{vertices})$

Rooted Trees (T^*) & Planar Rooted Trees (T)

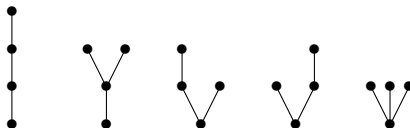
Definition

A **rooted n -tree** t is a connected finite graph with n vertices, no cycles, and a distinguished vertex r called the **root**. It is called **planar** if it is endowed with an ordering on vertices (to the children at each vertex).

- **parent of v :** unique vertex u adjacent to v , closer to the root.
- **children of u :** all vertices (if any) with u as parent.
- **leaf:** a vertex ℓ with no children.
- **subtree t_v rooted at v :** tree consisting of v and its descendants.
- **\mathbb{I} :** empty tree (size 0)



(a) rooted tree $t \in T^*$



(b) class T_4 of planar rooted 4-trees $t \in T$

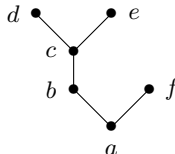
Antichain in the Set $V(t)$ of Vertices

Definition

An **antichain** in the set $V(t)$ of vertices of a rooted tree t is any subset $C \subseteq V(t)$ consisting of vertices that cannot be compared.

i.e., no two vertices in C have an ancestor-descendant relationship.

Examples: siblings *or* the set of all children in the same generation.



Antichains: $\{d, e, f\}$

$\{d, e\}, \{b, f\}, \{c, f\}, \{d, f\}, \{e, f\}$

$\{a\}, \{b\}, \{c\}, \{d\}, \{e\}, \{f\}$

\emptyset

Algebra of Planar Rooted Trees $k[T]$

k : field of characteristic 0

Definition

The **algebra of planar rooted trees** is the polynomial algebra $k[T]$ generated by the elements of T , with formal addition as addition and disjoint union as multiplication. The empty tree is identified with $1 \in k$.

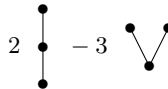
Definition

forest (n, m) **of size** $n = \sum_{i=1}^m n_i$ **(vertices) with** m **trees:**

Multiplication of m n_i -trees ($i = 1, \dots, m$).



(a) forest $(5, 2)$



(b) element in $k[T]$

Definition

Given a combinatorial class C , the **augmented generating function** of C is the formal power series

$$C(x) = \sum_{c \in C} cx^{|c|} \in (k[C])[[x]]$$

$$T^*(x) = \mathbb{I} + \bullet x + \begin{array}{c} \bullet \\ | \\ \bullet \end{array} x^2 + \left(\begin{array}{c} \bullet \\ | \\ \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet & \bullet \\ & \diagdown \diagup \\ & \bullet \end{array} \right) x^3 + \left(\begin{array}{c} \bullet \\ | \\ \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet & \bullet \\ & \diagdown \diagup \\ & \bullet \end{array} + \begin{array}{c} \bullet & \bullet & \bullet \\ & \diagdown \diagup & \\ & \bullet & \end{array} + \begin{array}{c} \bullet & \bullet & \bullet & \bullet \\ & \diagdown \diagup & \diagdown \diagup \\ & \bullet & \end{array} \right) x^4 \\ + O(x^5)$$

Denote: GF=generating function

AGF=augmented generating function

Definition

An **evaluation map** is a map $\phi : k[C] \longrightarrow \mathbb{Z}_{n \geq 0}$.

- **ordinary:** $\text{or}(c) = 1$
- **exponential:** $\text{ex}(c) = \frac{1}{|c|!}$

Note: defined on elements $c \in C$, extended to $k[C]$ as an algebra homomorphism.

- **ordinary GF:** $\text{or}(C(x)) = \sum_{c \in C} x^{|c|}$

$$\text{or}(T^*(x)) = 1 + x + x^2 + 2x^3 + 4x^4 + O(x^5) \quad (\text{coef}(x^n) = \# \text{ } n\text{-trees})$$

- **exponential GF:** $\text{ex}(C(x)) = \sum_{c \in C} \frac{x^{|c|}}{|c|!}$

$$\text{ex}(T^*(x)) = 1 + x + \frac{1}{2!}x^2 + \frac{2}{3!}x^3 + \frac{4}{4!}x^4 + O(x^5) \quad \left(\text{coef}(x^n) = \frac{\# \text{ } n\text{-trees}}{\# \text{ } n\text{-labels}} \right)$$

- Combinatorial objects
- Combine objects \longrightarrow product
- Decompose objects \longrightarrow coproduct
- “*inverse*” transformation \longrightarrow antipode

Connes-Kreimer Hopf Algebra of Rooted Trees

$H = k[T^*]$: the space of forests of rooted trees

\mathbb{I} : the empty forest (empty monomial)

Definition

The **Connes-Kreimer Hopf algebra of rooted trees** is the following combinatorial (commutative) Hopf algebra:

- **grading on H :** number of vertices of a forest
- **connected:** since $\text{span}(H_0) = k\mathbb{I} \cong k$
- **unit:** \mathbb{I}
- **counit:** $\varepsilon(t) = 0, \varepsilon(\mathbb{I}) = 1$
- **product:** disjoint union of rooted trees
- **coproduct:**
$$\Delta(t) = \sum_{\substack{C \subseteq V(t) \\ C \text{ antichain}}} \left(\prod_{v \in C} t_v \right) \otimes \left(t - \prod_{v \in C} t_v \right) \begin{pmatrix} \text{splitting:} \\ \text{"pruned"} \\ \text{\&} \\ \text{"stump"} \end{pmatrix}$$
- **antipode:**
$$S(t) = -t - \sum_{\substack{\emptyset \subsetneq C \subsetneq V(t) \\ C \text{ antichain}}} S\left(\prod_{v \in C} t_v\right) \left(t - \prod_{v \in C} t_v\right) \begin{pmatrix} \text{recursive} \\ \text{relation} \end{pmatrix}$$

- **Grading of forests:** measures how ecologically matured/grown/complex the forest is;

low grading: newly planted forest, young and sparse

high grading: thriving ecosystem of well-established, sprawling woodland

- **H connected:** empty forest has meaning; land that can be planted
- **Product:** plants trees one next to the other in a forest
- **Unit:** empty land
- **Counit:** probability P of absence of tree structure;

empty forest: $P(\mathbb{I}) = 1 = \varepsilon(\mathbb{I})$

any other tree: $P(t) = 0 = \varepsilon(t)$

answers the question: “*Can I plant here?*” 1 = YES, 0 = NO.

Hopf Structure Interpretation II

- Coproduct Δ :**
- encodes all different ways to decompose trees
 - cuts edges under vertices in antichain
 - results in a “pruned” part and the remaining “stump”



Figure: Antichains: \emptyset , $\{a\}$, $\{b\}$, $\{c\}$, $\{d\}$, $\{b, c\}$, $\{d, c\}$.

$$\Delta \left(\begin{array}{c} \bullet \\ | \\ \bullet \quad \bullet \\ / \quad \backslash \\ \bullet \end{array} \right) = \mathbb{I} \otimes_{\emptyset} \begin{array}{c} \bullet \\ | \\ \bullet \quad \bullet \\ / \quad \backslash \\ \bullet \end{array} + \begin{array}{c} \bullet \\ | \\ \bullet \quad \bullet \\ / \quad \backslash \\ \bullet \end{array} \otimes_a \mathbb{I} + \begin{array}{c} \bullet \\ | \\ \bullet \end{array} \otimes_b \begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \bullet \otimes_c \begin{array}{c} \bullet \\ | \\ \bullet \quad \bullet \\ / \quad \backslash \\ \bullet \end{array} + \bullet \otimes_d \begin{array}{c} \bullet \quad \bullet \\ / \quad \backslash \\ \bullet \end{array} + \begin{array}{c} \bullet \quad \bullet \\ | \quad | \\ \bullet \end{array} \otimes_{b,c} \bullet$$

$$+ \dots \otimes_{d,c} \begin{array}{c} \bullet \\ | \\ \bullet \end{array}$$

$$\Delta \left(\begin{array}{c} \bullet \\ | \\ \bullet \quad \bullet \end{array} \right) = \Delta \left(\begin{array}{c} \bullet \\ | \\ \bullet \end{array} \right) \Delta \left(\begin{array}{c} \bullet \end{array} \right) = \left(\mathbb{I} \otimes \begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ | \\ \bullet \end{array} \otimes \mathbb{I} + \bullet \otimes \bullet \right) \left(\mathbb{I} \otimes \bullet + \bullet \otimes \mathbb{I} \right)$$

$$= \mathbb{I} \otimes \begin{array}{c} \bullet \\ | \\ \bullet \quad \bullet \end{array} + \bullet \otimes \begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ | \\ \bullet \end{array} \otimes \bullet + \begin{array}{c} \bullet \quad \bullet \\ | \quad | \\ \bullet \end{array} \otimes \mathbb{I} + \bullet \otimes \bullet \bullet + \bullet \bullet \otimes \bullet$$

Hopf Structure Interpretation III

- Antipode S :**
- recursively dismantles trees; flips sign at each step
 - coproduct splits tree; antipode compensates for splitting

$$\begin{aligned}
 S\left(\begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \end{array}\right) &= -\begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \end{array} - \frac{1}{b} S\left(\begin{array}{c} \bullet \\ | \\ \bullet \end{array}\right) \begin{array}{c} \bullet \\ | \\ \bullet \end{array} - \frac{1}{c} S(\bullet) \begin{array}{c} \bullet \\ | \\ \bullet \end{array} - \frac{1}{d} S(\bullet) \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \end{array} - \frac{1}{b,c} S\left(\begin{array}{c} \bullet \quad \bullet \\ | \quad | \\ \bullet \end{array}\right) \bullet \\
 &\quad - \frac{1}{d,c} S(\bullet \bullet) \begin{array}{c} \bullet \\ | \\ \bullet \end{array} \\
 &= -\begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \end{array} + \begin{array}{c} \bullet \quad \bullet \\ | \quad | \\ \bullet \end{array} - 2 \begin{array}{c} \bullet \quad \bullet \\ | \quad | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ | \\ \bullet \end{array} \begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \end{array} + \bullet \bullet \bullet
 \end{aligned}$$

$$S(\bullet) = -\bullet$$

$$S\left(\begin{array}{c} \bullet \\ | \\ \bullet \end{array}\right) = -\begin{array}{c} \bullet \\ | \\ \bullet \end{array} - S(\bullet) \bullet = -\begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \bullet \bullet$$

$$S(\bullet \bullet) = -\bullet \bullet - 2 S(\bullet) \bullet = -\bullet \bullet + 2 \bullet \bullet = \bullet \bullet$$

$$\begin{aligned}
 S\left(\begin{array}{c} \bullet \quad \bullet \\ | \quad | \\ \bullet \end{array}\right) &= -\begin{array}{c} \bullet \quad \bullet \\ | \quad | \\ \bullet \end{array} - S(\bullet) \bullet \bullet - S(\bullet) \begin{array}{c} \bullet \\ | \\ \bullet \end{array} - S(\bullet \bullet) \bullet - S(\bullet \bullet) \bullet \\
 &= -\begin{array}{c} \bullet \quad \bullet \\ | \quad | \\ \bullet \end{array} + \bullet \bullet \bullet + \begin{array}{c} \bullet \\ | \\ \bullet \end{array} \bullet - 2 \bullet \bullet \bullet = -\bullet \bullet \bullet
 \end{aligned}$$

scattering amplitude A = probability of this outcome ($|A|^2$)

particles sent in \rightarrow collide & interact \rightarrow detect **what** comes out

What happened?

NO IDEA

perturbative expansion = take **weighted sum** over all **possibilities**

Particular scenario: graph in spacetime

edges = propagating particles & *vertices* = interactions

forget spacetime embedding
(*when & where*)

combine all possibilities
reduced to same abstract graph

1 possibility = 1 Feynman Graph
(*which* particles interacted & *how*)

weight = Feynman Integral

Cake interpretation of previous diagram:

- scattering amplitude = finished cake
- perturbative expansion = recipe
- sum = mixing ingredients & cooking
- graphs in spacetime = packaged ingredients (we know label & position)
- forgetting spacetime embedding = unpacking the ingredients
- possibilities = ingredients (= Feynman graphs)
- weight = quantity of ingredient (= Feynman Integral)

Feynman Integrals:

- **Coupling Constants:**

- ★ capture strength of each interaction
- ★ contribute to Feynman integral
- ★ can be interpreted as counting variables

- **Feynman Rules:**

- ★ how to read Feynman integrand expression off the graph
- ★ each edge & each vertex contribute a factor

- **Renormalization:**

- ★ extracting physically meaningful quantities from divergent integrals

- **Sums of Feynman Integrals:**

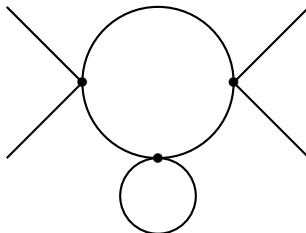
- ★ expected to be divergent for all interesting cases.

Definition

A **half-edge Feynman graph** G is a set H of half edges along with

- a set $V(G)$ of **vertices**: vertex=disjoint subset of half edges
- a set $E(G)$ of **internal edges**: internal edge=disjoint pair of half edges.

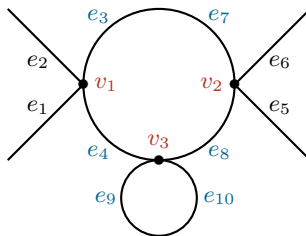
External edges=half edges not in internal edges.



Definition

A **half-edge labelled Feynman graph** G_L is a half-edge graph along with

- a labelling L : bijection $H \longrightarrow \{1, 2, \dots, |H|\}$.



$$V(G) = \{v_1, v_2, v_3\}, \quad E(G) = \{i_1, i_2, i_3\}$$

$$i_1 = \{e_3, e_7\}, \quad i_2 = \{e_4, e_8\}, \quad i_3 = \{e_9, e_{10}\}$$

$$v_1 = \{e_1, e_2, e_3, e_4\}, \quad v_2 = \{e_5, e_6, e_7, e_8\}, \quad v_3 = \{e_4, e_8, e_9, e_{10}\}$$

Theorem

Let G be a connected graph with n half edges. Let m be the number of half-edge labelled graphs giving G upon forgetting the labelling. Then:

$$\frac{m}{n!} = \frac{1}{|\text{Aut}(G)|} \quad (\text{symmetry factor})$$

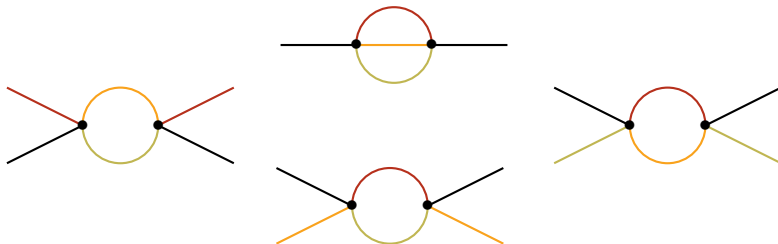
Why do we care?

- $L(x) = \text{aug}(\text{labelled})$ & $U(x) = \text{aug}(\text{unlabelled})$
- $\phi = \text{“forgetful” evaluation map, } \phi(\text{labelled}) = \phi(\text{unlabelled})$
- $\phi_L(\cdot) = \frac{1}{n!} \phi(\cdot)$ & $\phi_U(\cdot) = \frac{1}{|\text{Aut}(\cdot)|} \phi(\cdot)$
- $\phi_L(L(x)) = \phi_U(U(x))$
- $\text{or } (\text{unlabelled; weighted}) = \text{ex}(\text{labelled})$
- **FACT:** $\text{ex}(\text{labelled}) = \exp(\text{ex}(\text{connected}))$
- $\mathbf{A}_{\Sigma \text{ all}} = \exp(\mathbf{A}_{\Sigma \text{ connected}})$

Definition

One-particle irreducible or **1PI** Feynman graph:

- connected
- connected after removing any one internal edge.



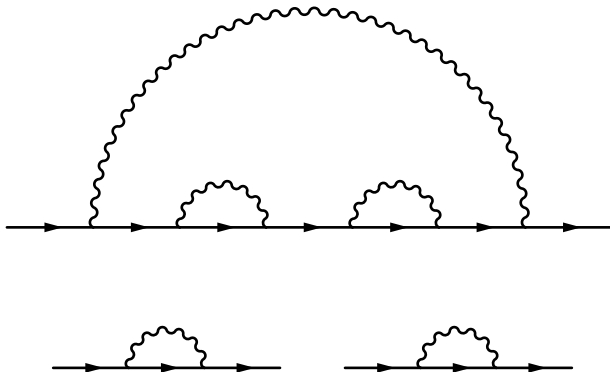
$$\mathbf{1PI} \xleftarrow[\text{Transform}]{\text{Legendre}} \text{connected graphs} \xrightarrow{\text{exp}} \text{graphs}$$

Analogy:

- connected graphs = walls
- 1PI graphs = bricks
- non-1PI graphs = part of wall with multiple bricks

Definition

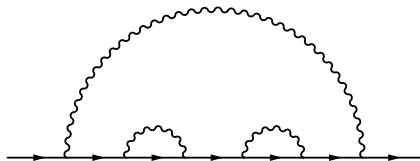
subgraph γ of graph G : subset of vertices of G



Divergent & Primitive Feynman Graphs

Definition

A Feynman graph is **divergent** if it has a divergent Feynman integral. A divergent Feynman graph with no divergent subgraphs is called **primitive**.



Definition

A subgraph γ is called **subdivergence** if its connected components are 1PI and divergent.

Overlapping subdivergences: they share at least 1 vertex.



Definition

A **combinatorial physical theory** T is a set of half edge types with

- **permissible edge types:** set of types of edges (pairs of half edges)
- **permissible vertex types:** set of types of vertices
- **power counting weight** w : an integer for each type (edge & vertex)
- **dimension** D of spacetime: a nonnegative integer

Note: edges here are *internal* !

Definition

superficial degree of divergence of G in T :

$$\omega = D\ell - \sum_e w(e) - \sum_v w(v)$$

- $\omega \geq 0$: G **divergent**
- $\omega = 0$: G **logarithmically divergent**

Definition

external leg structure $L(G)$: multiset of half edge types of external edges of G

Definition

A combinatorial physical theory T is **renormalizable** if:

$$\omega(G, T) = f(L(G)), \quad \forall G \in T$$

$$\phi^k : \begin{cases} 1 \text{ half-edge type} \\ 1 \text{ edge type, } w(e) = 2 \\ 1 \text{ } k\text{-valent vertex, } w(v) = 0 \end{cases} \quad \begin{array}{l} \bullet \phi^4 : D = 4 \\ \bullet \phi^3 : D = 6 \end{array}$$

Denote: $\ell = \#$ loops, $v = \#$ vertices, $e = \#$ edges, $x = \#$ external legs

Euler characteristic: $\ell = e - v + 1$

Regularity equation: $kv = x + 2e$

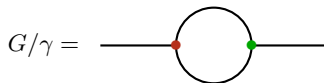
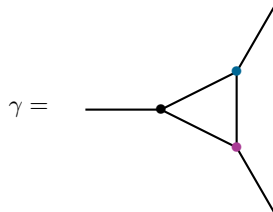
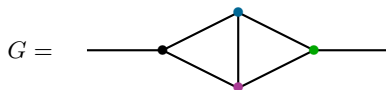
$$\begin{aligned} \text{Then: } \omega &= D\ell - \sum_e w(e) - \sum_v w(v) \\ &= D\ell - 2e \\ &= D + \left(1 - \frac{D}{2}\right)x + \underbrace{\left(\frac{Dk}{2} - D - k\right)v} \\ \implies D &= \frac{2k}{k-2} \end{aligned}$$

Definition

Let G be a Feynman graph in a theory T . Let γ be a subgraph with each connected component 1PI and divergent. The **contraction** G/γ of γ is the Feynman graph in T constructed as follows:

- start from G
- identify subgraph γ
- collapse vertices and internal edges into a single vertex
- keep external edges attached to external vertices
- this is G/γ .

Contraction of γ : Example



Definition

Fix a renormalizable theory T . Let \mathcal{G} be the set of connected 1PI graphs in T . The **renormalization Hopf algebra** \mathcal{H} associated to T is the polynomial algebra $k[\mathcal{G}]$ with:

- **generators:** divergent 1PI graphs of T
- **grading:** loop order
- **unit:** empty graph \mathbb{I}
- **counit:** $\epsilon(\mathbb{I}) = 1$, $\epsilon(G) = 0$ if G nonempty
- **product:** disjoint union of graphs
- **coproduct:**
$$\Delta(G) = \sum_{\substack{\gamma \subseteq G \\ \gamma = \gamma_p}} \gamma \otimes G/\gamma$$
- **antipode:**
$$S(G) = -G - \sum_{\substack{\emptyset \subsetneq \gamma \subsetneq G \\ \gamma = \gamma_p}} S(\gamma) G/\gamma$$

γ_p : product of divergent 1PI subgraphs

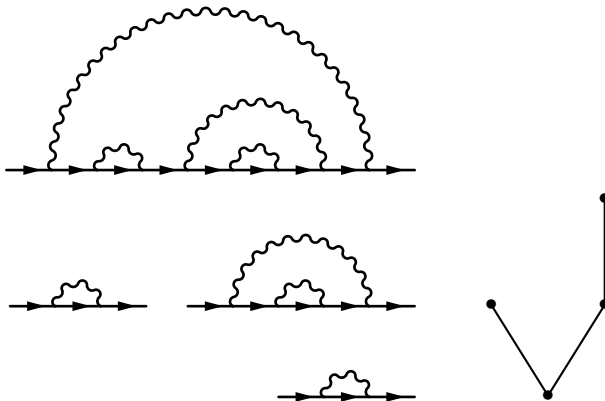
Coproduct: Example (in ϕ^3)

$$\begin{aligned}
 & \Delta \left(\text{---} \text{---} \text{---} \right) \\
 &= \mathbb{I} \otimes \text{---} \text{---} \text{---} + \text{---} \text{---} \text{---} \otimes \mathbb{I} \\
 &+ 2 \text{---} \text{---} \text{---} \otimes \text{---} \text{---} \text{---}
 \end{aligned}$$

The diagrams represent Feynman graphs in the ϕ^3 theory. The first diagram is a diamond shape with a vertical line through its center, representing a self-energy correction. The second and third diagrams show the coproduct of this diagram with the identity element \mathbb{I} on either side. The fourth diagram shows a triangle with three external lines, representing a 1-loop vertex correction, which is multiplied by 2 and then tensored with a circle with two external lines, representing a 1-loop self-energy correction.

Note: γ with 1 or 2 vertices is NOT 1PI, so not taken into account!

Insertion Trees: Rooted Trees Take The Stage!



The End *...or maybe Not*

The slides are full of mathematical *Easter Eggs* ! Have Fun !!

- [1] Yeats, K. *A Combinatorial Perspective on Quantum Field Theory*. Springer, 2017.
- [2] Cartier, P., and Patras, F. *Classical Hopf Algebras and Their Applications*. Springer, 2021.
- [3] Panzer, E. *Hopf-algebraic Renormalization of Kreimer's Toy Model*. Master's thesis, Humboldt University of Berlin, 2012.

Appendix A: Feynman Diagrams in Mathematics & Physics

- Algebraic Geometry:
 - Moduli Spaces
 - Intersection Theory
- Geometric Topology:
 - Knot Invariants
 - 3-Manifolds
- Representation Theory:
 - of Quantum Groups
 - of Lie Algebras
- Number Theory:
 - Prime Numbers
 - Riemann Zeta Function
- Probability and Statistics:
 - Stochastic Processes
 - Random Walks
- Combinatorics:
 - Graphical Representations
 - Generating & Partition Functions
- Graph Theory:
 - Connectivity, cycles, paths
 - Graph Algorithms
- Mathematical Physics:
 - Statistical Mechanics
 - Perturbative Expansions
- Quantum Computing:
 - Quantum Circuits
 - Quantum Algorithms
- Artificial Intelligence:
 - Neural Networks
 - AI Systems/Models

$$\begin{aligned}\omega &= D\ell - \sum_e w(e) - \sum_v w(v) \\ &= D\ell - 2e \\ &= D(e - v + 1) - kv + x \\ &= D\left(\frac{kv}{2} - \frac{x}{2} - v + 1\right) - kv + x \\ &= D + \left(\frac{Dk}{2} - D - k\right)v + \left(1 - \frac{D}{2}\right)x\end{aligned}$$

Euler characteristic: $\ell = e - v + 1$

regularity equation: $kv = x + 2e$